Centrality Measures on Big Graphs:
Exact, Approximated, and Distributed Algorithms

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ABSTRACT

Centrality measures allow to measure the relative impor-
tance of a node or an edge in a graph w.r.t. other nodes or
edges. Several measures of centrality have been developed
in the literature to capture different aspects of the informal
concept of importance, and algorithms for these different
measures have been proposed. In this tutorial, we survey
the different definitions of centrality measures and the algo-
rithms to compute them. We start from the most common
measures, such as closeness centrality and betweenness cen-
trality, and move to more complex ones such as spanning-
edge centrality. In our presentation, we begin from exact al-
gorithms and then progress to approximation algorithms, in-
cluding sampling-based ones, and to highly-scalable MapRe-
duce algorithms for huge graphs, both for exact computation
and for keeping the measures up-to-date on dynamic graphs
where edges are inserted or removed over time. Our goal is
to show how advanced algorithmic techniques and scalable
systems can be used to obtain efficient algorithms for an im-
portant graph mining task, and to encourage research in the
area by highlighting open problems and possible directions.

Keywords

Centrality; Betweenness; Closeness; Tutorial

1. INTRODUCTION

Identifying the "important" nodes or edges in a graph is
a fundamental task in network analysis, with many appli-
cations, from economics and biology to security and soci-
ology. Several measures, known as *centrality indices*, have
been proposed over the years, formalizing the concept of
importance in different ways [16]. Centrality measures rely
on graph properties to quantify importance. For example,
betweenness centrality, one of the most commonly used cen-
trality indices, counts the fraction of shortest paths going
through a node, while the closeness centrality of a node

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is the average sum of the inverse of the distance to other
nodes. Other centrality measures use eigenvectors, random
walks, degrees, or more complex properties. For instance,
the PageRank index of a node is a centrality measure, and
centrality measures for sets of nodes have also been defined.

With the proliferation of huge networks with millions of
nodes and billions of edges, the importance of having scal-
able algorithms for computing centrality indices has become
more and more evident, and a number of contributions have
been recently proposed, ranging from heuristics that per-
form extremely well in practice to approximation algorithms
offering strong probabilistic guarantees, to scalable algo-
rithms for the MapReduce platform. Moreover, the dynamic
nature of many networks, i.e., the addition and removal of
nodes or edges over time, dictates the need to keep the com-
puted values of centrality up-to-date as the graph changes.
These challenging problems have enjoyed enormous interest
from the research community, with many relevant contribu-
tions proposed recently to tackle them.

Our tutorial presents, in a unified framework, some of the
many measures of centrality, and discusses the algorithms to
compute them, both in an exact and in an approximate way,
both in-memory and in a distributed fashion in MapReduce.
The goal of this unified presentation is to ease the compar-
ison between different measures of centrality, the different
quality guarantees offered by approximation algorithms, and
the different trade-offs and scalability behaviors characteriz-
ing distributed algorithms. We believe this unity of presen-
tation is beneficial both for newcomers and for experienced
researchers in the field, who will be exposes to the material
from a coherent point of view.

The graph analyst can now choose among a huge number
of centrality indices, from the well-established ones origi-
nally developed in sociology, to the ones more recently in-
troduced ones that capture other aspects of importance. At
the same time, the original algorithms that could handle
the relatively small networks for classic social science exper-
iments have been superseded by important algorithmic con-
tributions that exploit modern computational frameworks
or obtain fast, high-quality approximations. It is our be-
lief that the long history of centrality measures and the
ever-increasing interest from computer scientists in analyz-
ing larger and richer graphs create the need for an all-around
organization of both old and new materials, and the desire
to satisfy this need inspired us to develop this tutorial.

2. OUTLINE

The tutorial is structured in three main technical parts, plus a concluding part where we discuss future research directions. All three technical parts will contains both theory and experimental results.

Part I: Definitions and Exact Algorithms.

In this first part, we introduce the different centrality measures, starting from important axioms that a good centrality measure should require. We then discuss the relationship between the different measures, including results highlighting the high correlation between many of them. After having laid these foundations, we move to present the algorithms for the exact computation of centrality measures, both in static and in dynamic graph. We discuss the state-of-the-art by presenting both algorithms with the best worst-case time complexity and heuristics that work extremely well in practice by exploiting different properties of real world graphs.

Exact computation of centrality measures becomes impractical on web-scale networks. Commonly, one of two alternative approaches is taken to speed up the computation: focus on obtaining an *approximation* of the measure of interest or use *parallel and distributed algorithms*. In the second part of our tutorial we explore the former approach, while in the third part we deal with the latter.

Part II: Approximation Algorithms.

Most approximation algorithms for centrality measures use various forms of sampling and more or less sophisticated analysis to derive a sample size sufficient to achieve the desired level of approximation with the desired level of confidence. In this part we present a number of these sampling based algorithms, from simple ones using the Hoeffding inequality to more complex ones using VC-dimension and Rademacher Averages. For each algorithm, we discuss its merits and drawbacks, and highlight the challenges for the algorithm designer. We also discuss the case of maintaining an approximation up-to-date in a dynamic graphs, presenting a number of contributions that recently appeared in the literature.

Part III: Highly-scalable Algorithms.

In the third part of our tutorial, we discuss parallel and distributed algorithms for the computation of centrality measures in static and dynamic graphs. Specifically we present an approach based on GPUs and one based on processing parallel/distributed data streams, together with experimental results.

List of topics with references.

The following is a preliminary list of topics we will cover in each part of the tutorial, with the respective references.

1. **Introduction: definitions and exact algorithms**

   (a) The axioms of centrality [3]

   (b) Definitions of centrality [16], including, but not limited to: betweenness, closeness, degree, eigenvector, harmonic, Katz, absorbing random-walk [14], and spanning-edge centrality [13].

   (c) Betweenness centrality: exact algorithm [4] and heuristically-faster exact algorithms for betweenness centrality [7, 22].

   (d) Exact algorithms for betweenness centrality in a dynamic graph [12, 15, 17].

   (e) Exact algorithms for closeness centrality in a dynamic graph [21].

2. **Approximation algorithms**

   (a) Sampling-based algorithm for closeness centrality [6].

   (b) Betweenness centrality: almost-linear-time approximation algorithm [23], basic sampling-based algorithm [5], refined estimators [8], VC-dimension bounds for betweenness centrality [18], Rademacher bounds for betweenness centrality [19].

   (c) Approximation algorithms for betweenness centrality in dynamic graphs [1, 2, 9, 10].

3. **Highly-scalable algorithms**

   (a) GPU-based algorithms [20].

   (b) Exact parallel streaming algorithm for betweenness centrality in a dynamic graph [11].

4. **Challenges and directions for future research**

3. INTENDED AUDIENCE

The tutorial is aimed at researchers interested in the theory and the applications of algorithms for graph mining and social network analysis.

We do not require any specific existing knowledge. The tutorial is designed for an audience of computer scientists who have a general idea of the problems and challenges in graph analysis. We will present the material in such a way that any advanced undergraduate student would be able to productively follow our tutorial and we will actively engage with the audience and adapt our pace and style to ensure that every attendee can benefit from our tutorial. The tutorial starts from the basic definitions and progressively moves to more advanced algorithms, including sampling-based approximation algorithms and MapReduce algorithms, so that it will be of interest both to researchers new to the field and to a more experienced audience.

4. SUPPORT MATERIALS

We developed a mini-website for the tutorial at `http://matteo.rionda.to/centrtutorial/`. It contains the abstract of the tutorial, a detailed outline with short a description of each item of the outline, a full list of references with links to electronic editions, a list of software packages implementing the algorithms, and the slides used in the tutorial presentation.

5. INSTRUCTORS

This tutorial is developed by Francesco Bonchi, Gianmarco De Francisci Morales, and Matteo Riondato. All three instructors will attend the conference.

Francesco Bonchi is Research Leader at the ISI Foundation, Turin, Italy, where he leads the "Algorithmic Data Analytics" group. He is also Scientific Director for Data Mining at Eurecat (Technological Center of Catalunya), Barcelona. Before he was Director of Research at Yahoo Labs in Barcelona, Spain, leading the Web Mining Research group.

His recent research interests include mining query-logs, social networks, and social media, as well as the privacy issues related to mining these kinds of sensible data.

He will be PC Chair of the 16th IEEE International Conference on Data Mining (ICDM 2016) to be held in Barcelona in December 2016. He is member of the ECML PKDD Steering Committee, Associate Editor of the newly created IEEE Transactions on Big Data (TBD), of the IEEE Transactions on Knowledge and Data Engineering (TKDE), the ACM Transactions on Intelligent Systems and Technology (TIST), Knowledge and Information Systems (KAIS), and member of the Editorial Board of Data Mining and Knowledge Discovery (DMKD). He is co-editor of the book "Privacy-Aware Knowledge Discovery: Novel Applications and New Techniques" published by Chapman & Hall/CRC Press.

He presented a tutorial at ACM KDD'14.

Gianmarco De Francisci Morales is a Scientist at QCRI. Previously he worked as a Visiting Scientist at Aalto University in Helsinki, as a Research Scientist at Yahoo Labs in Barcelona, and as a Research Associate at ISTI-CNR in Pisa. He received his Ph.D. in Computer Science and Engineering from the IMT Institute for Advanced Studies of Lucca in 2012. His research focuses on scalable data mining, with an emphasis on Web mining and data-intensive scalable computing systems. He is an active member of the open source community of the Apache Software Foundation, working on the Hadoop ecosystem, and a committer for the Apache Pig project. He is one of the lead developers of Apache SAMOA, an open-source platform for mining big data streams. He commonly serves on the PC of several major conferences in the area of data mining, including WSDM, KDD, CIKM, and WWW. He co-organizes the workshop series on Social News on the Web (SNOW), co-located with the WWW conference. He presented a tutorial on stream mining at IEEE BigData'14.

Matteo Riondato is a Research Scientist in the Labs group at Two Sigma Investments. Previously he was a postdoc at Stanford and at Brown. His dissertation on sampling-based randomized algorithms for data and graph mining received the Best Student Poster Award at SIAM SDM'14. His research focuses on exploiting advanced theory in practical algorithms for time series analysis, pattern mining, and social network analysis. He presented tutorials at ACM KDD'15, ECML PKDD'15, and ACM CIKM'15.

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