# **Specification of Web Services**

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Transparency of business processes is critical for cooperation

#### Needs

semantic constraints temporal aspects access control policies consider non-web processes

## Current Approaches

can't model relations b/w inputs and outputs differentiate b/w outputs and effects mostly miss / no semantics of ACPs don't support

### Our Approach

Combination of Pi-Calculus and DLs formal semantics for temporal aspects variables to model relations b/w inputs and outputs access control policies as pre-conditions

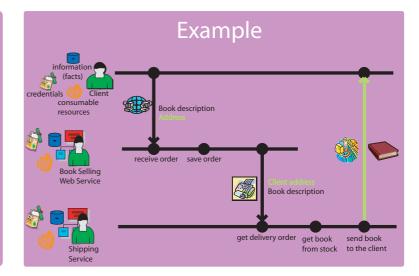
#### **Features**

Unified specification of outputs and effects Channel = (protocol, address, message type) can model non-web processes as well captures correct semantics of comm. Interoperable credential based access control

## Syntax

$$A \coloneqq 0 \mid$$
 null  $l(v_1,...,v_j) \mid$  input activity  $\bar{l}o_1,...,o_k \mid$  output activity  $[c]P,Q$  if-then-else

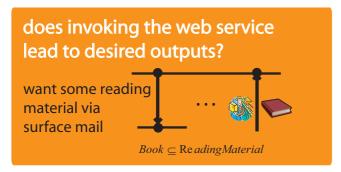
$$W := A_1 \parallel \ldots \parallel A_n$$
 composition



# **Automatic Matchmaking of Web Services**

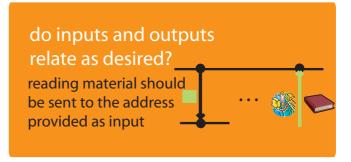
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expressive request specification formalism

### Matchmaking Algorithm

Given a request R and a web service description W calculate V(R), the set of variables in R calculate V(W), the set of variables in W from V(R) and V(W) calculate the substitution functions  $\sigma_1, \dots, \sigma_n$  For each i in  $\{1, \dots, n\}$  generate  $W_i$  from W and  $\sigma_i$  by  $\alpha$ -conversion eliminate condition expressions from  $W_i \to W_{i,1}, \dots, W_{i,2}$   $C_{i,1}, \dots, C_{i,2}$  For each j in  $\{1, \dots, 2^{k+1}\}$  If the structures of R and  $W_{i,j}$  are same and every input type of R is sub concept of corr. input type of  $W_{i,j}$  and every output rule of R subsumes corr. output rule of  $W_{i,j}$  and every condition in R is subsumed by the corr. condition of  $W_{i,j}$  and then W is a match for R with conditions  $C_{i,j}$