



# Probabilistic Automated Bidding in Alternative Auctions

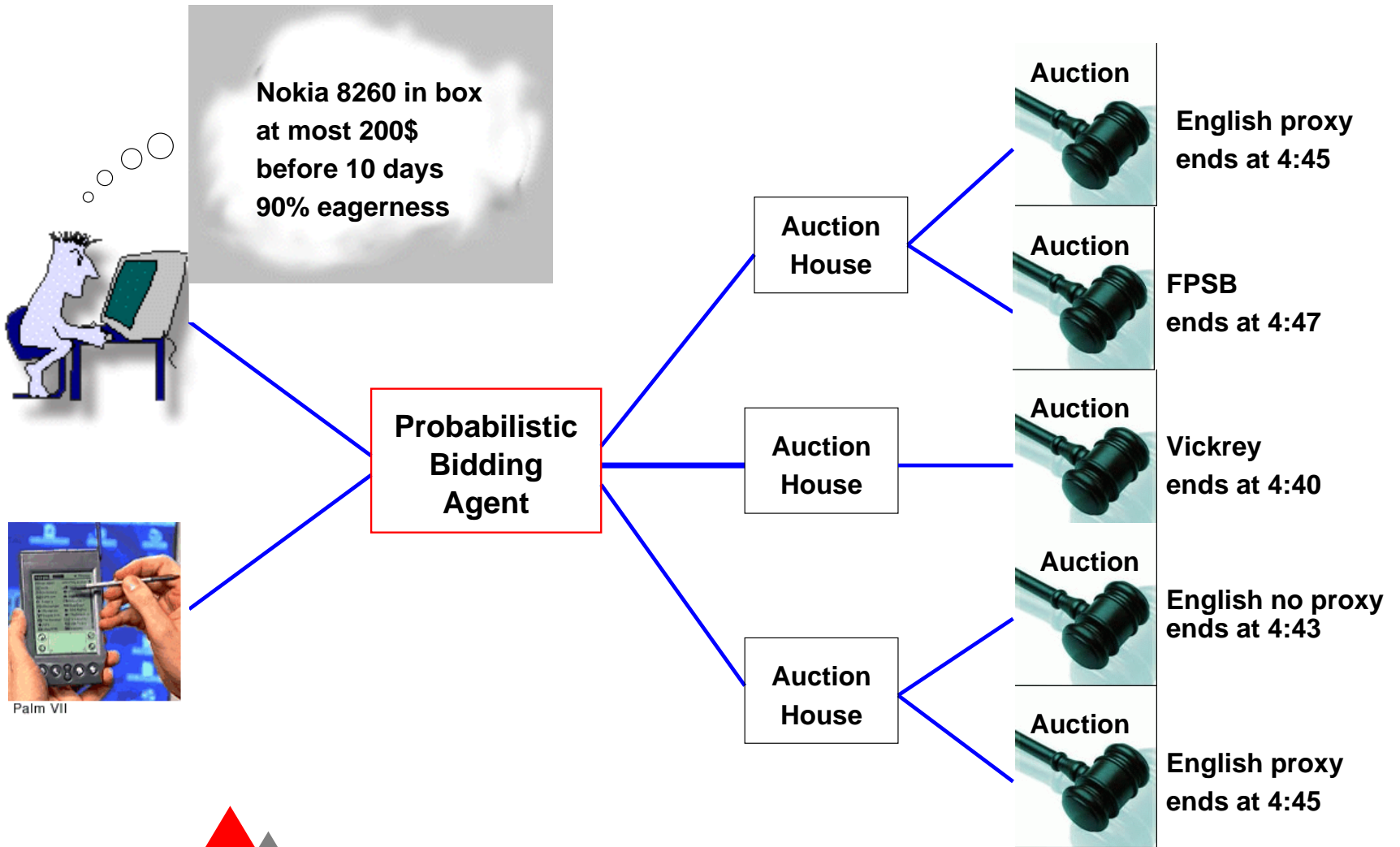
**Marlon Dumas**, Lachlan Aldred, Guido  
Governatori, Arthur ter Hofstede, Nick Russell

Queensland University of Technology, Australia

**[m.dumas@qut.edu.au](mailto:m.dumas@qut.edu.au)**



# Vision





# Goal

To obtain one unit of an item at the lowest price, given the following parameters:

$M$ : The maximum bidding price

$D$ : The deadline for obtaining the item

$G$ : The *eagerness* to obtain the item



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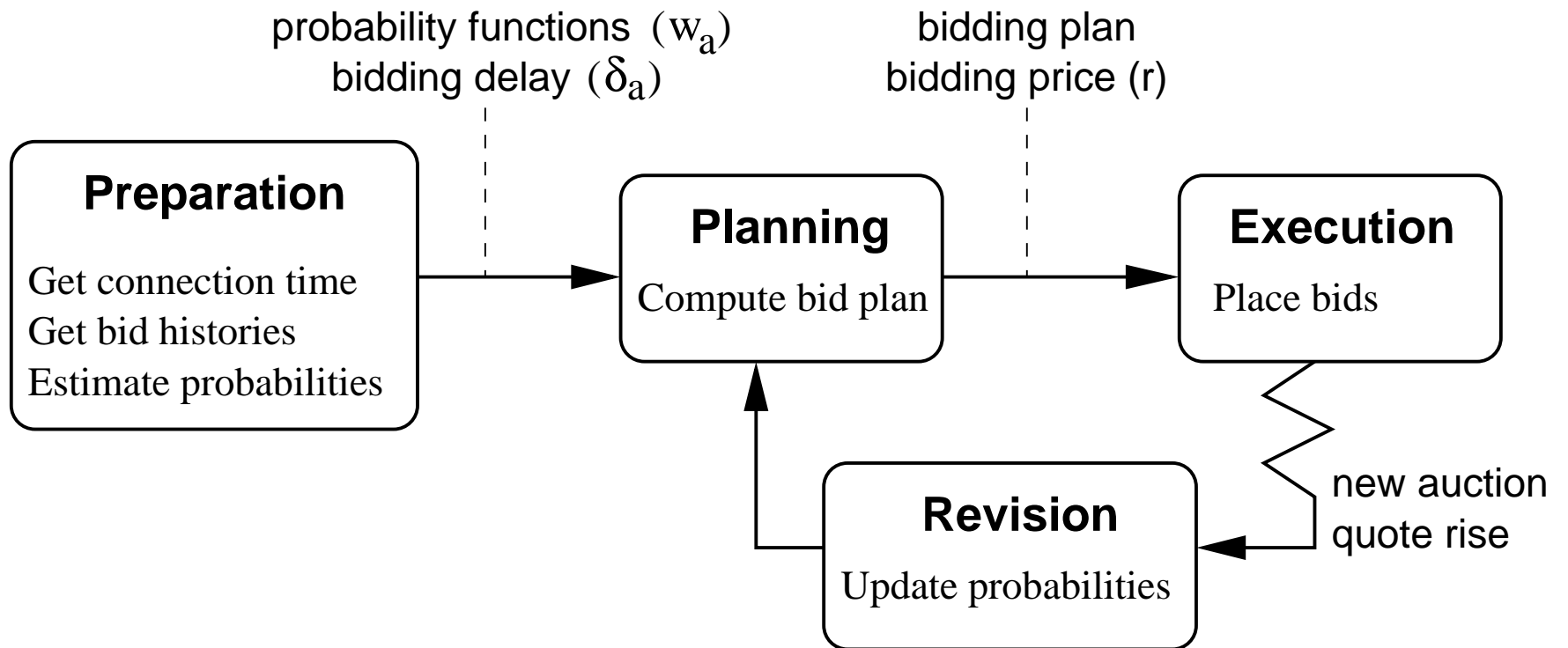
$G$ : The *eagerness* to obtain the item

Auctions are single-unit with fixed deadlines:

- eBay-style auctions with or without proxy bids
- FPSB and Vickrey auctions

# Approach

A bidding agent operates in 4 phases:





## Preparation: Probability estimation

Given the history of Winning Bids (W.B.) and the quote  $q$  of an auction, the probability of winning with a bid of  $r$  can be computed in two ways.

### Histogram method

$$w(r) = \frac{\# \text{ of auctions with W.B. between } q \text{ and } r}{\# \text{ of auctions with W.B. greater than } q}$$





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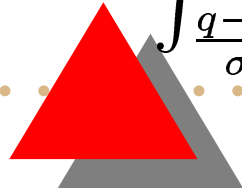
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### Normal distribution method

$$w(r) = \frac{\int_{\frac{q-\mu}{\sigma}}^{\frac{r-\mu}{\sigma}} e^{-x^2/2} dx}{\int_{\frac{q-\mu}{\sigma}}^{\infty} e^{-x^2/2} dx}$$

$\mu$  = average W.B.  
 $\sigma$  = std. dev. of W.B.

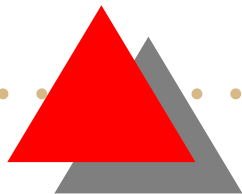




# Planning: Problem statement

Given a set  $A_a$  of announced auctions, find:

- A set of auctions  $A_s \subseteq A_a$
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such that:

- Auctions in  $A_s$  are mutually compatible

$$\forall a_1, a_2 \in A_s \quad |\text{end}(a_2) - \text{end}(a_1)| \geq \delta_{a_1} + \delta_{a_2}$$

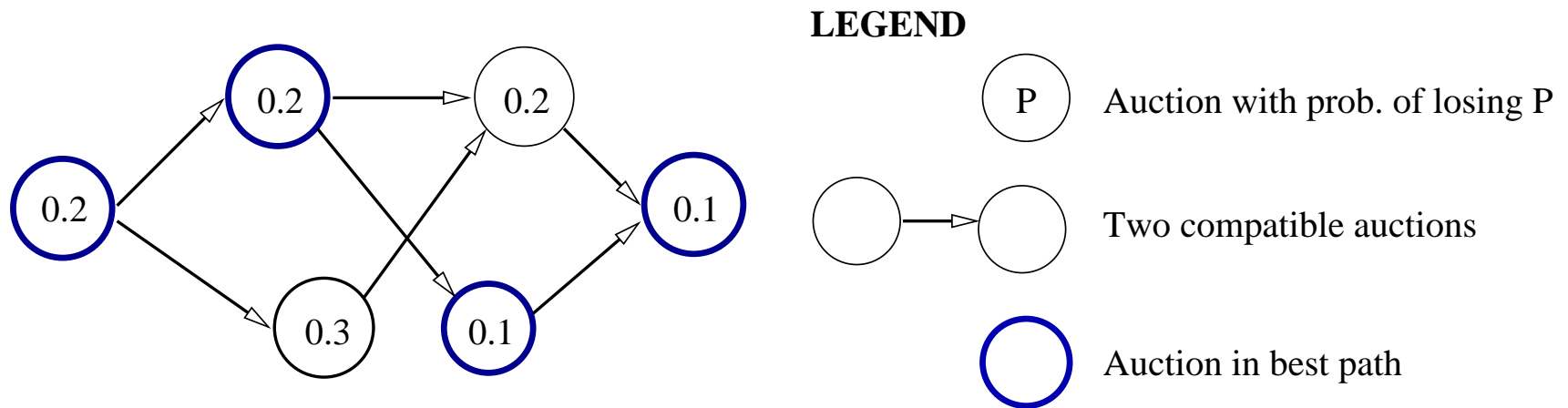
- Probability of winning 1 auction is satisfactory

$$1 - \prod_{a \in A_s} (1 - w_a(r)) \geq G$$

- $r$  is minimal w.r.t. the previous constraints

# Planning: Computing the best plan

For a given price  $r$ , it is possible to compute the best bidding plan using a *critical path algorithm*.



Prob. of loosing in best plan =  $.2^2 \times .1^2 = .004$

Prob. of winning in best plan =  $1 - .004 = 99.6\%$



# Planning: Minimising the bidding price

For each  $r$  between 1 and  $M$

    Compute the best bidding plan at price  $r$  ;

    If the prob. of winning with this plan is  $\geq G$ ,  
    stop iterating

If no appropriate  $r$  is found, notify the user.

Otherwise, take  $r$  as the bidding price.

**Note:** Binary search can be used as optimisation





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A plan revision is triggered in the following cases:

- A new auction for the required item appears
- The quote of an auction in the plan rises above the bidding price



# Heterogeneity between auctions

*Alternative auctions* are often heterogeneous:

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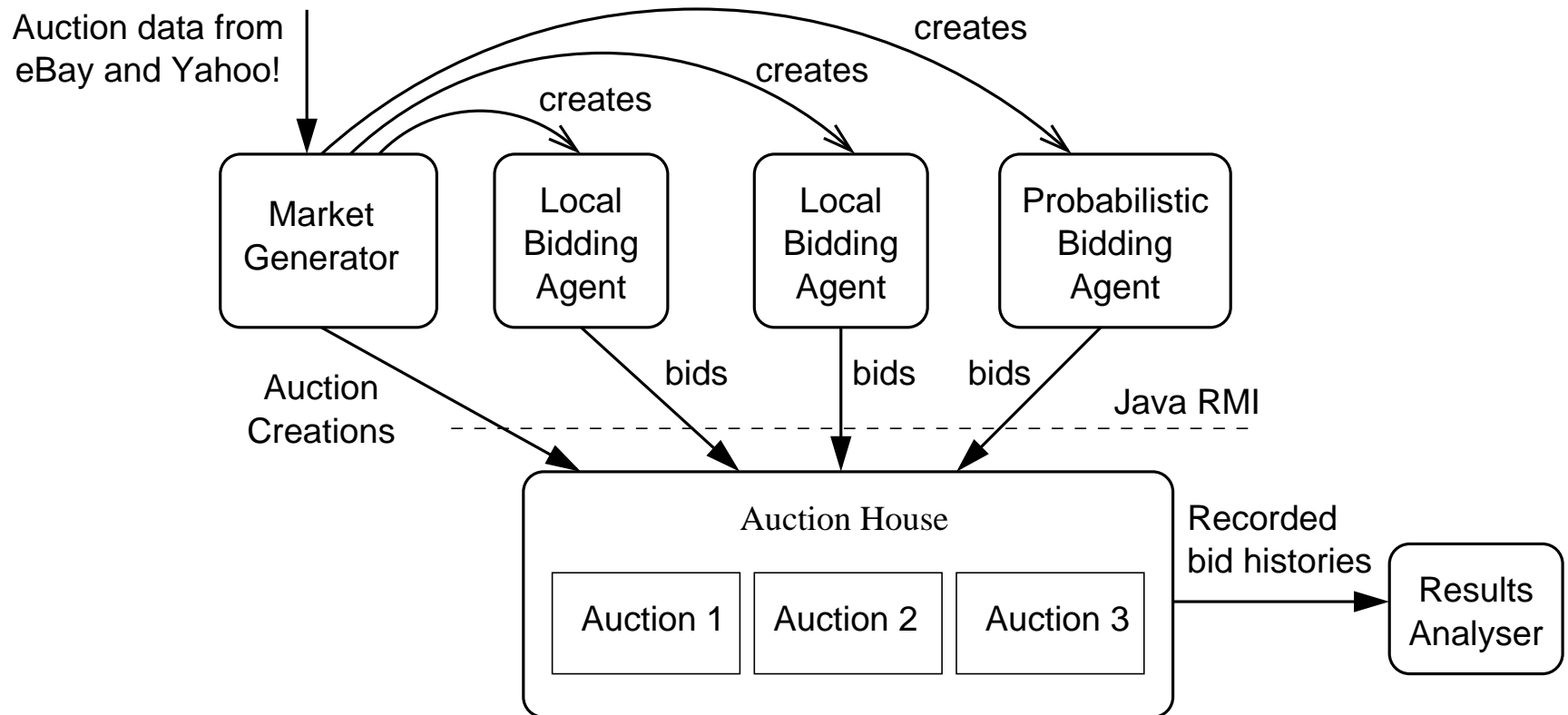
Two approaches to deal with heterogeneity:

- Price differentiation. The user sets a different maximum price for each auction
- Utility differentiation. The user provides a multi-attribute scoring system



# Experimentation

## Auction simulation platform





# Experimentation

## Tested claims

1. The percentage of times that a probabilistic bidder wins is equal to its eagerness

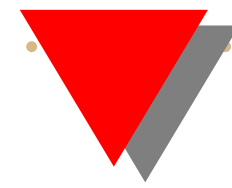


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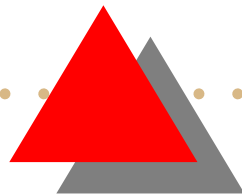
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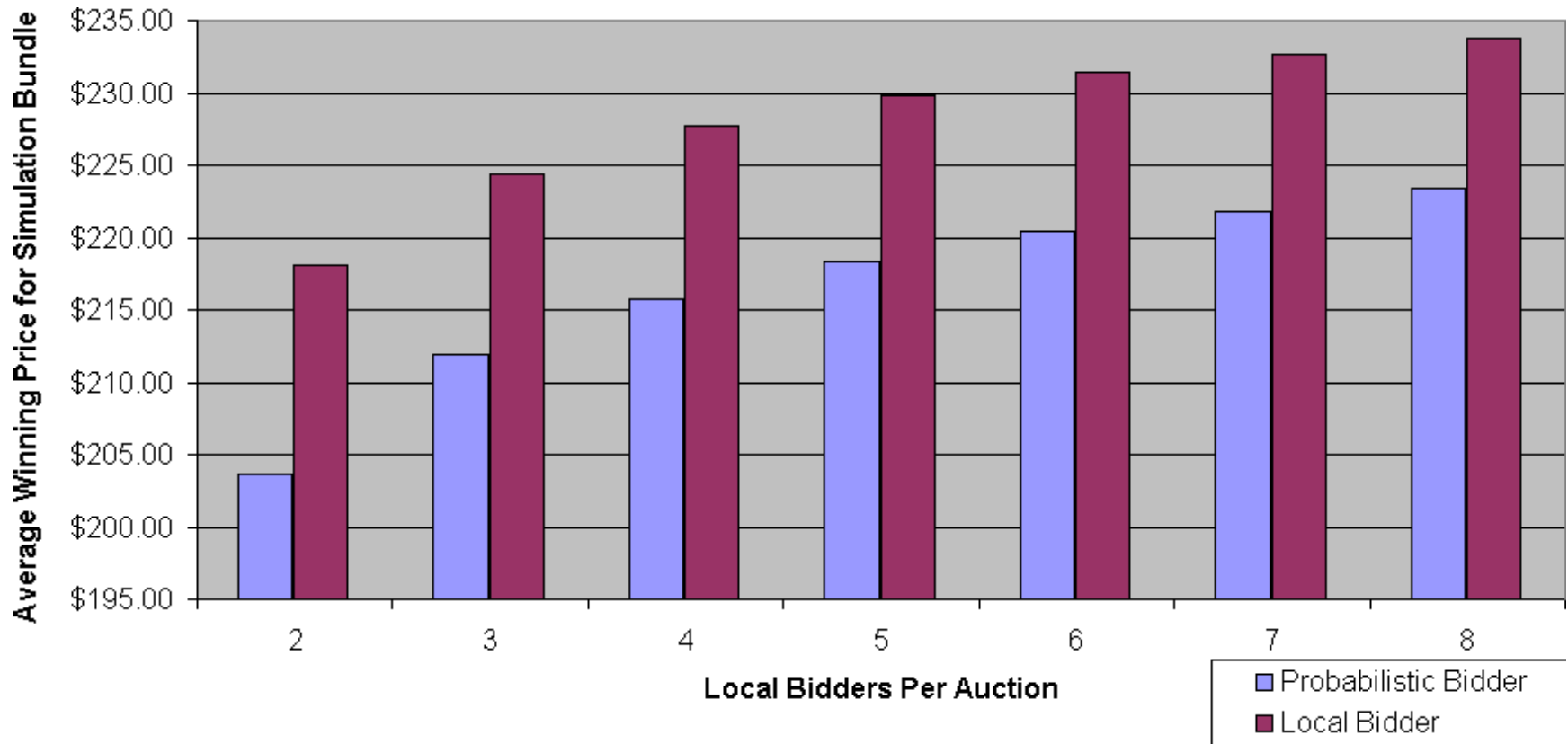
## Tested claims

1. The percentage of times that a probabilistic bidder wins is equal to its eagerness
2. Probabilistic bidders pay less than local ones
3. The welfare of the market increases with the number of probabilistic bidders



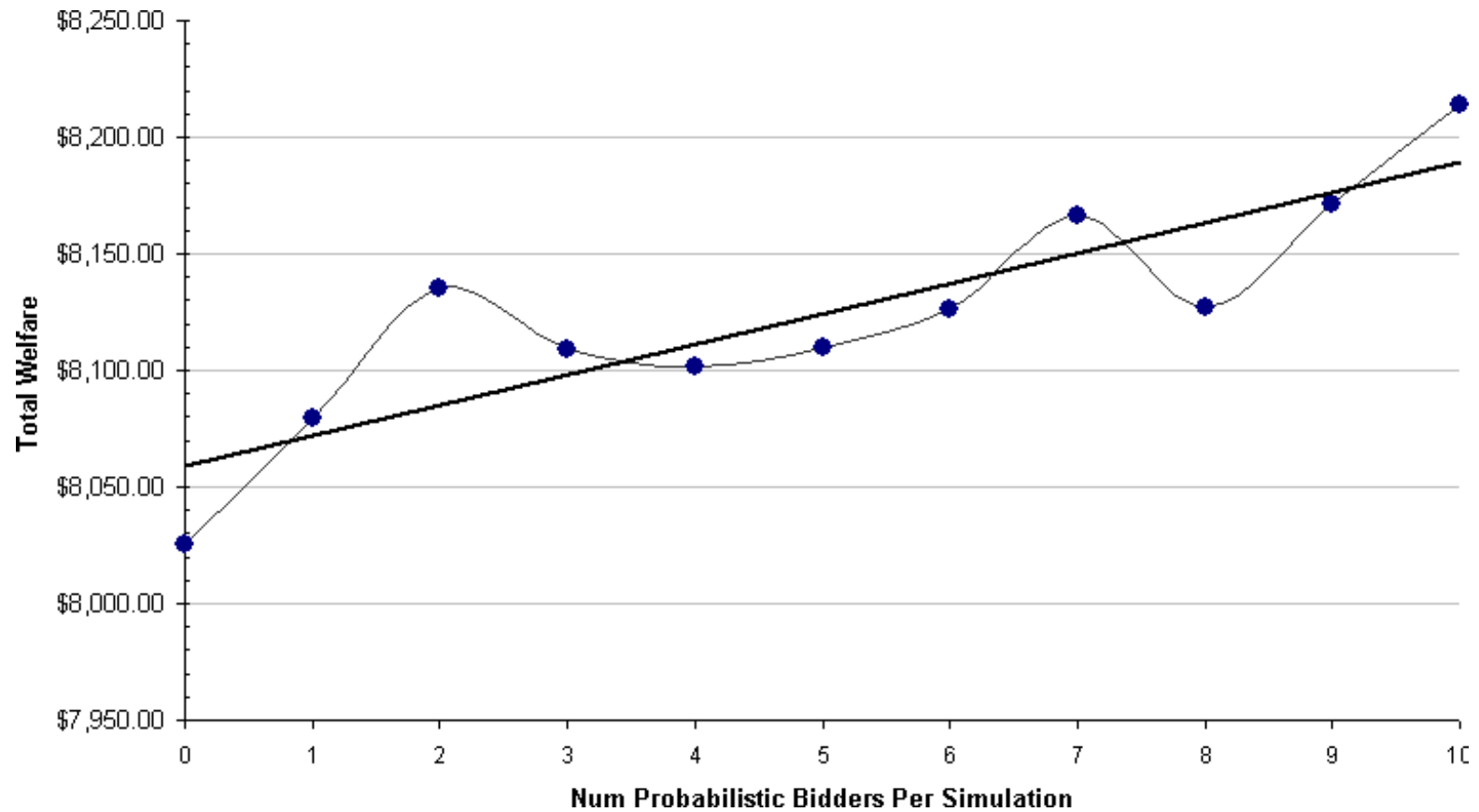
# Experimentation

## Validation of Claim 2



# Experimentation

## Validation of Claim 3





# Conclusion

## Probabilistic bidding agents:

- allow bidders to make tradeoffs between price and eagerness;
- increase the payoff of their users and the welfare of the market

## Future extensions:

- Multiple units of an item / multi-unit auctions
- Interrelated items (all-or-none transactions)

