

A lightweight protocol for the generation and distribution of secure e-coupons

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Outline of the talk

- Motivations
- A PKI based solution
- The lightweight protocol
 - The static model
 - The dynamic model
 - The extended model

The motivations

- The web and the market
 - E-commerce
 - Advertising
- Advertising on the web
 - Web servers host ads
 - Non-active role
 - Banners, product info, etc.
 - Web servers distribute promotional material
 - Active role
 - Discounts, gifts, etc.

From coupons to e-coupons

- Paper coupons



-traditional form of advertisement distributed by

- mail
- fliers
- newspapers and magazines
- ...



-redemption at physical store



- Hybrid coupons

-downloaded from the web or received by e-mail



-printed out



-redemption at physical store



E-coupons

- Truly digital

- download from advertisers' web sites
- redemption at stores' web site



- Advantages



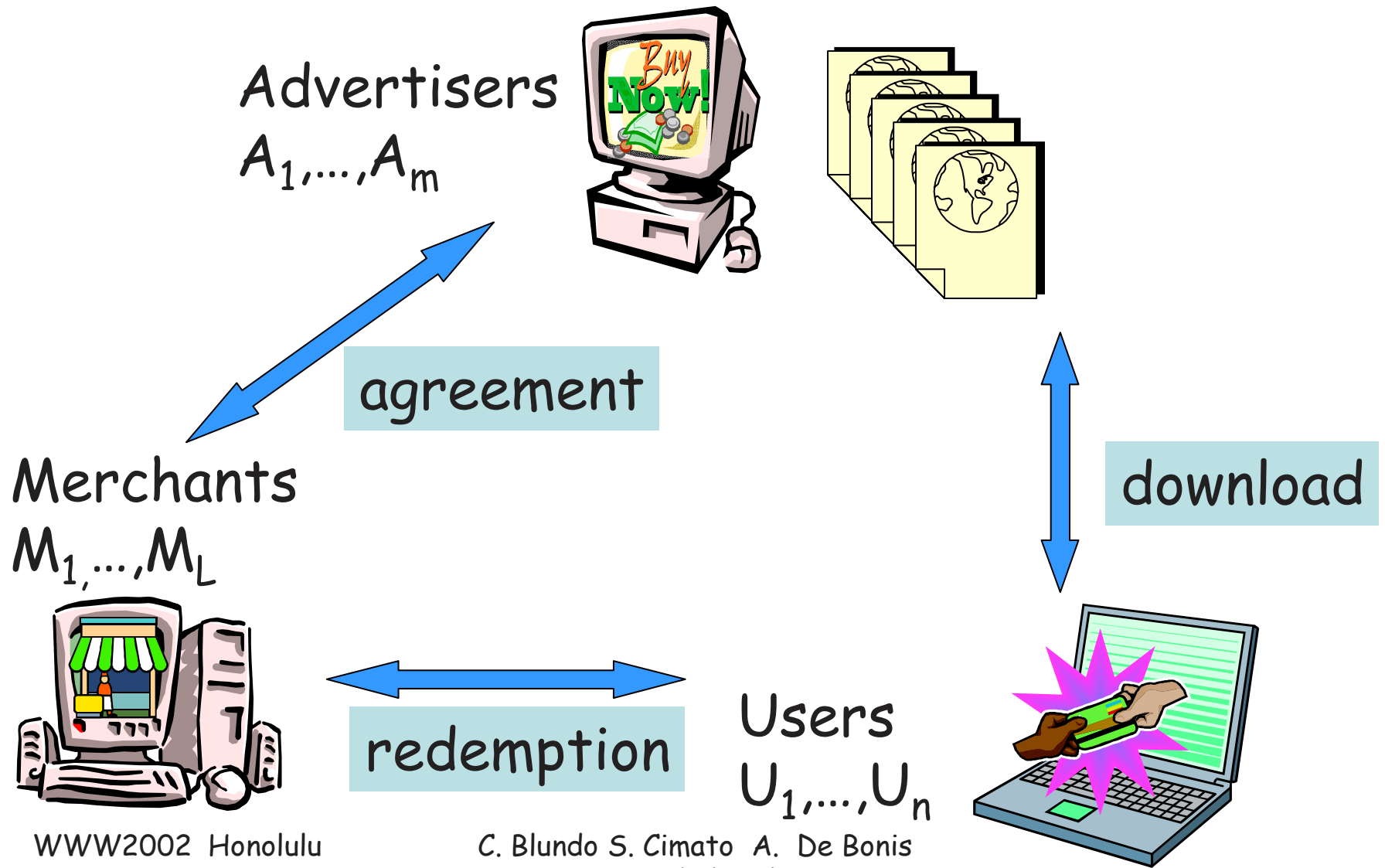
- ✓ Easy to distribute, to store and to use
- ✓ Target recipients
- ✓ Feedback on promotional campaign

- Disadvantages

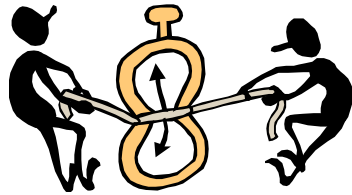
- ✓ Duplication
- ✓ Double spending
- ✓ Forging
- ✓ Manipulation



The scenario & life cycle



Requirements



Practical Requirements

- Soundness
- Efficiency
- Ease to use
- Interoperability
- Customers' anonymity

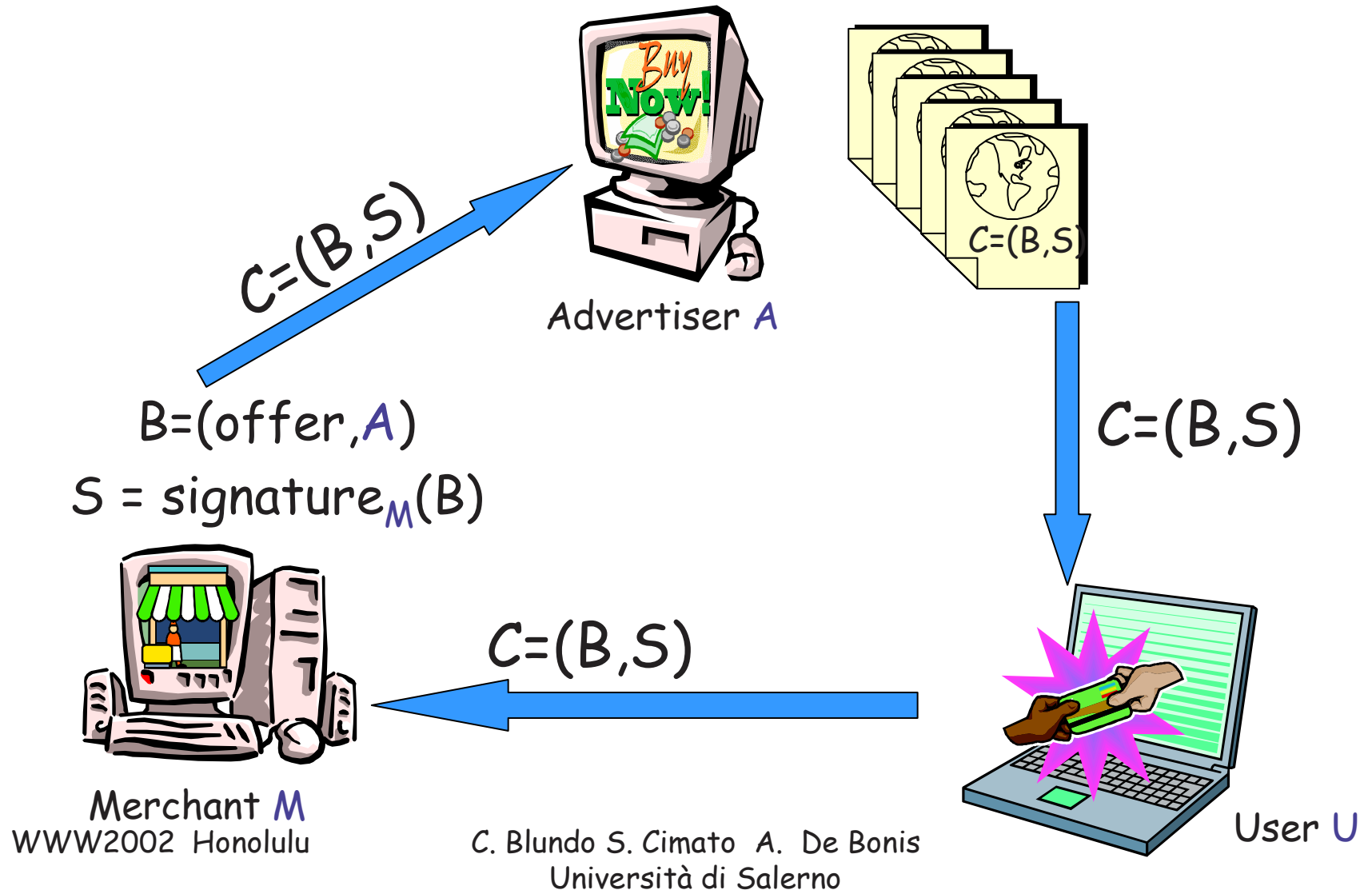


Security Requirements

- No unauthorized issuance
- No manipulation
- No double spending

A solution based on digital signature

[Jakobsson et al., WWW8]



Merchant M
WWW2002 Honolulu

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User U

Analysis of Jakobsson et al. solution

- Security Requirements:

Yes No unauthorized issuance

Yes No manipulation

Yes Double spending (if record of each e-coupon is maintained)

- Practical Requirements:

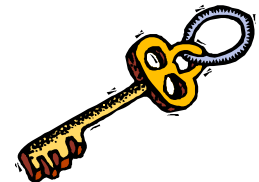
Yes Soundness

Yes Ease to use

Yes Interoperability

Yes Customers' anonymity

No Efficiency : Public key operations are very costly



A lightweight protocol



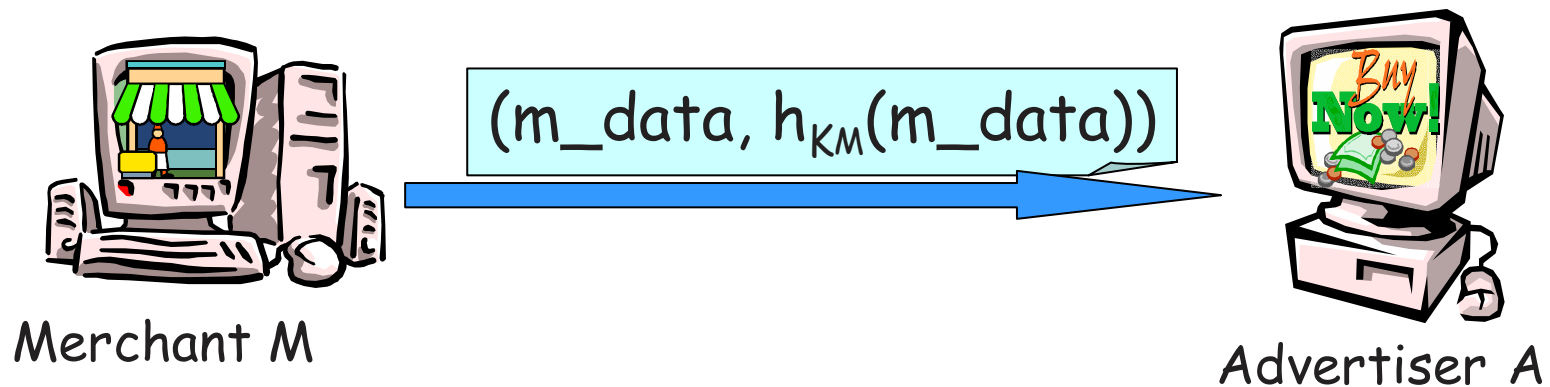
Idea: Use Mac Algorithms instead of digital signatures

Mac Algorithm:

- encoding functions $h_k(\cdot)$ parameterized by secret key k
 - easy to compute
 - compression: finite bit length output string
 - computational resistance: if k unknown, it is computationally infeasible to compute $h_k(x)$

Initialization phase

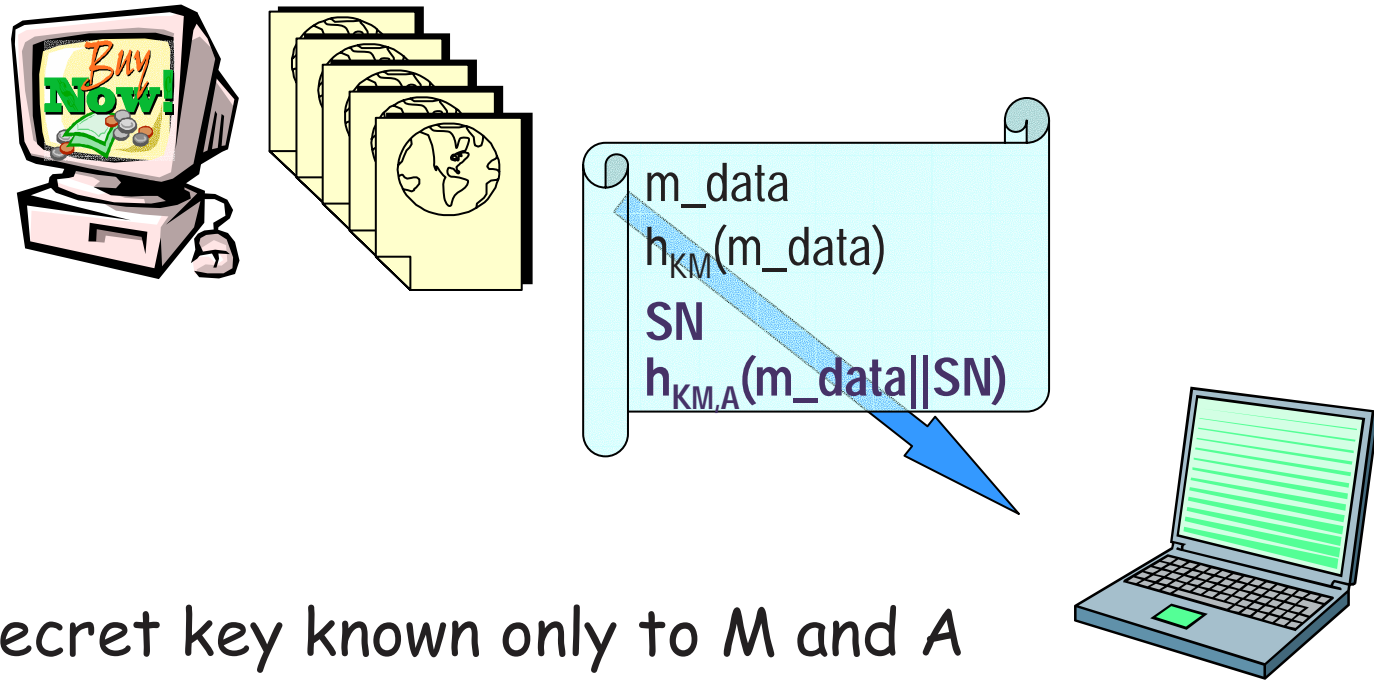
- Merchant M sends the e-coupon **framework** to advertiser A



- k_M : secret key known only to M
- h_{k_M} : MAC algorithm with parameter k_M
- m_data : merchant id, advertiser id, offer info,...

Generation phase

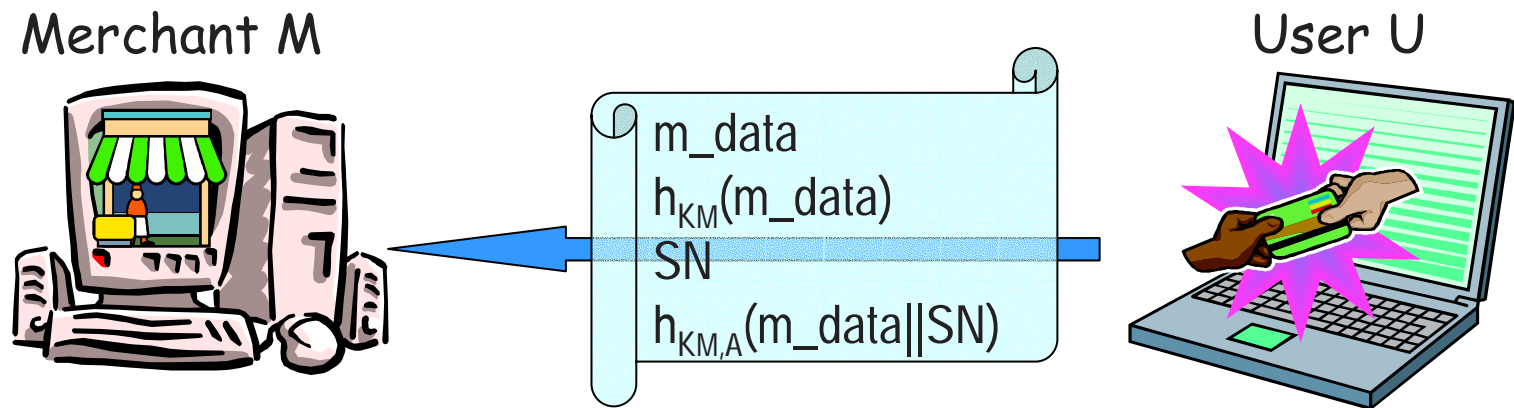
Advertiser A generates e-coupons using M's framework



- $k_{M,A}$: secret key known only to M and A
- h_{KMA} : MAC algorithm with parameter $k_{M,A}$
- SN : e-coupon serial number

Redemption phase

User U sends the e-coupon to merchant M



M computes $h_{KM}(m_data)$ and $h_{KM,A}(m_data||SN)$
and accepts e-coupon **iff**

- computed values = values stored in e-coupon
- SN has not been seen before

Security

- No unauthorized issuance:
 - Users cannot compute $h_{KM}(m_data)$ and $h_{KM,A}(m_data||SN)$
 - Advertiser cannot compute $h_{KM}(m_data||SN)$
- No manipulation:
 - Changing e-coupon data requires to compute new values of $h_{KM}(m_data)$ and $h_{KM,A}(m_data||SN)$
- No double-spending:
 - No two accepted e-coupons with the same Serial Number

Dynamic e-coupons

- Dynamic e-coupons contain information on the release time:
 - *Aging e-coupon*: offer decreases from the moment e-coupon is downloaded
- Dynamic e-coupons are released for marketing purposes:
 - Discourage users from shopping around for better offers
 - Give immediate feedback on advertisement campaign

Dynamic e-coupons



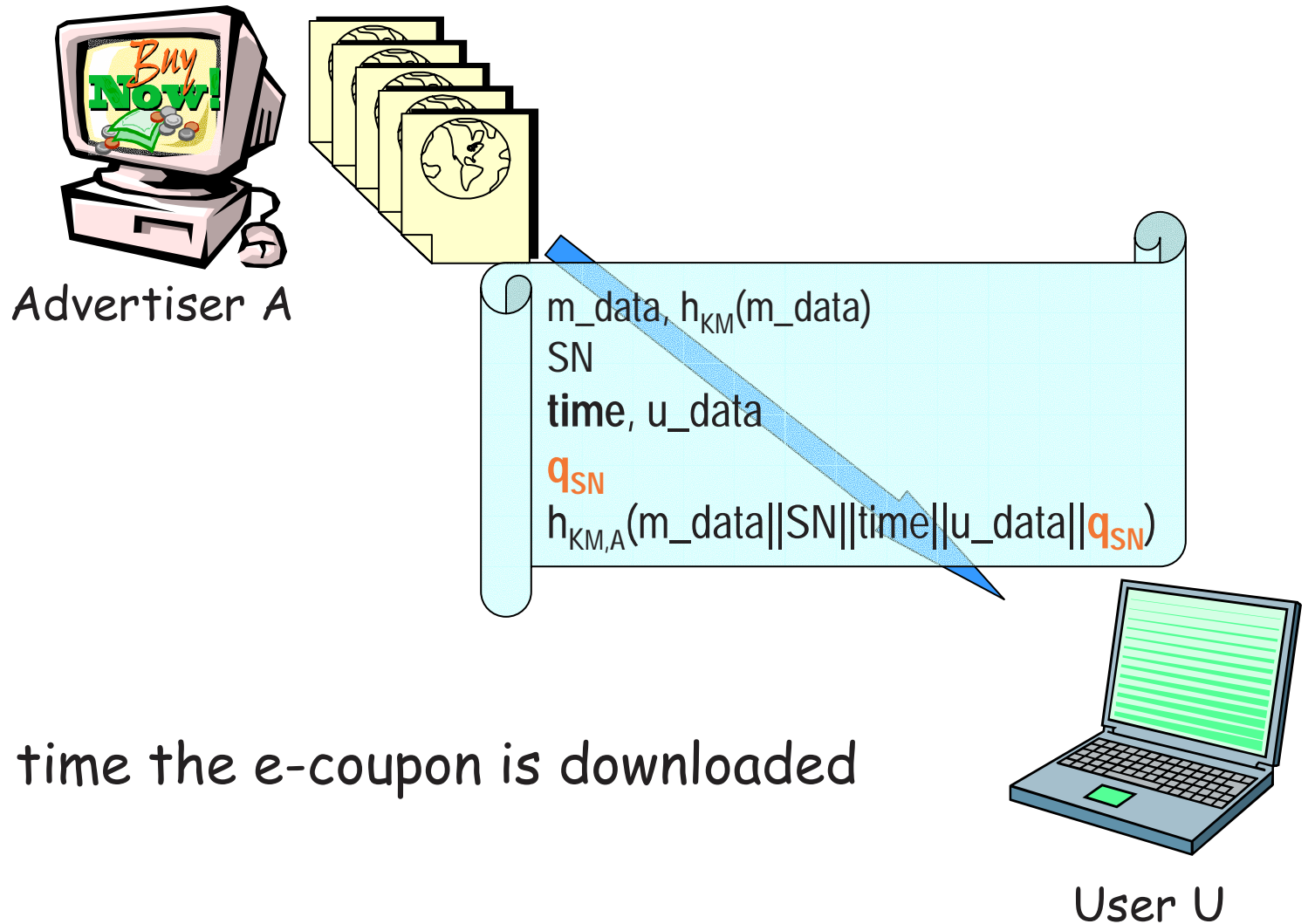
Idea: released e-coupons form a dependence chain

$q_1, q_2, q_3, \dots,$



- Advertiser computes:
 - chain "ring" for released e-coupon
 - $q_{SN} = q(u_data || q_{SN-1})$ (SN = serial number)
- Needed:
 - u_data : info released by user (e.g., IP address)
 - q : **collision resistant** hash function
 - It is computationally infeasible to find distinct x and y such that $q(x) = q(y)$

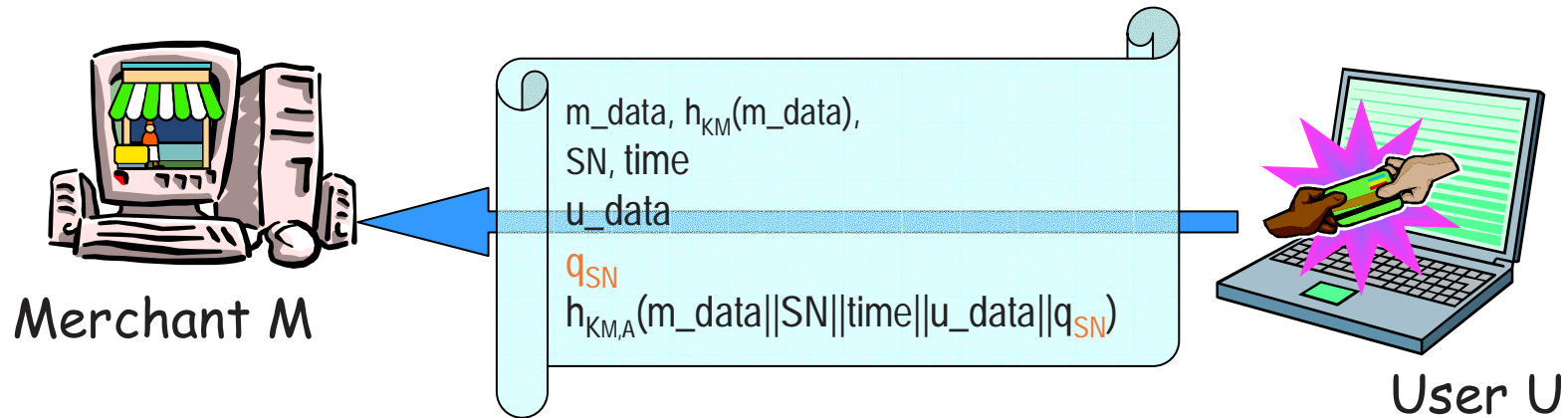
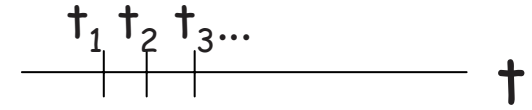
Generation of Dynamic E-coupons



- $time$: time the e-coupon is downloaded

Redemption phase

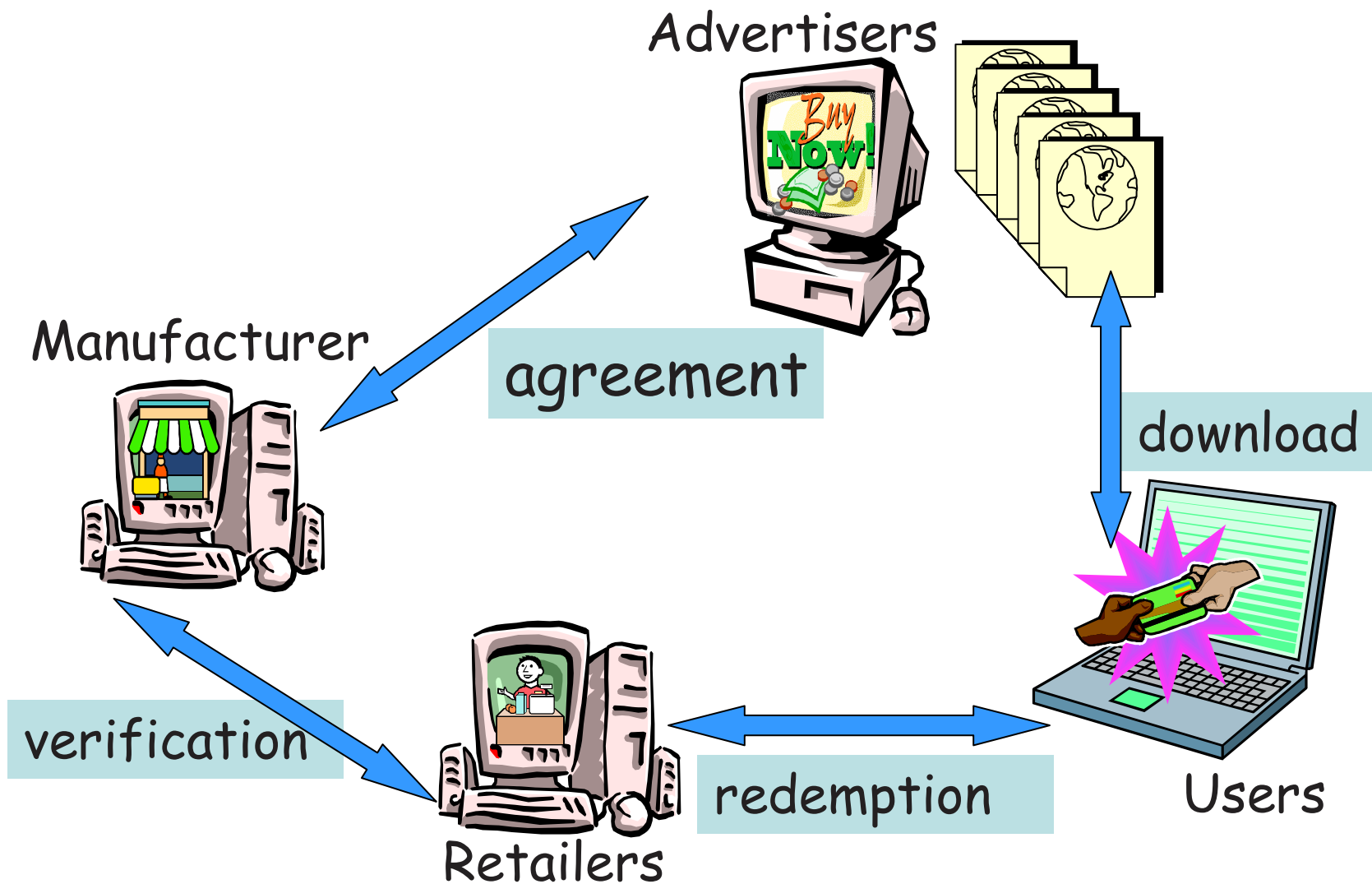
- Time is divided in time frames t_1, t_2, t_3, \dots
- A gives M the dependence chains for previous time frames



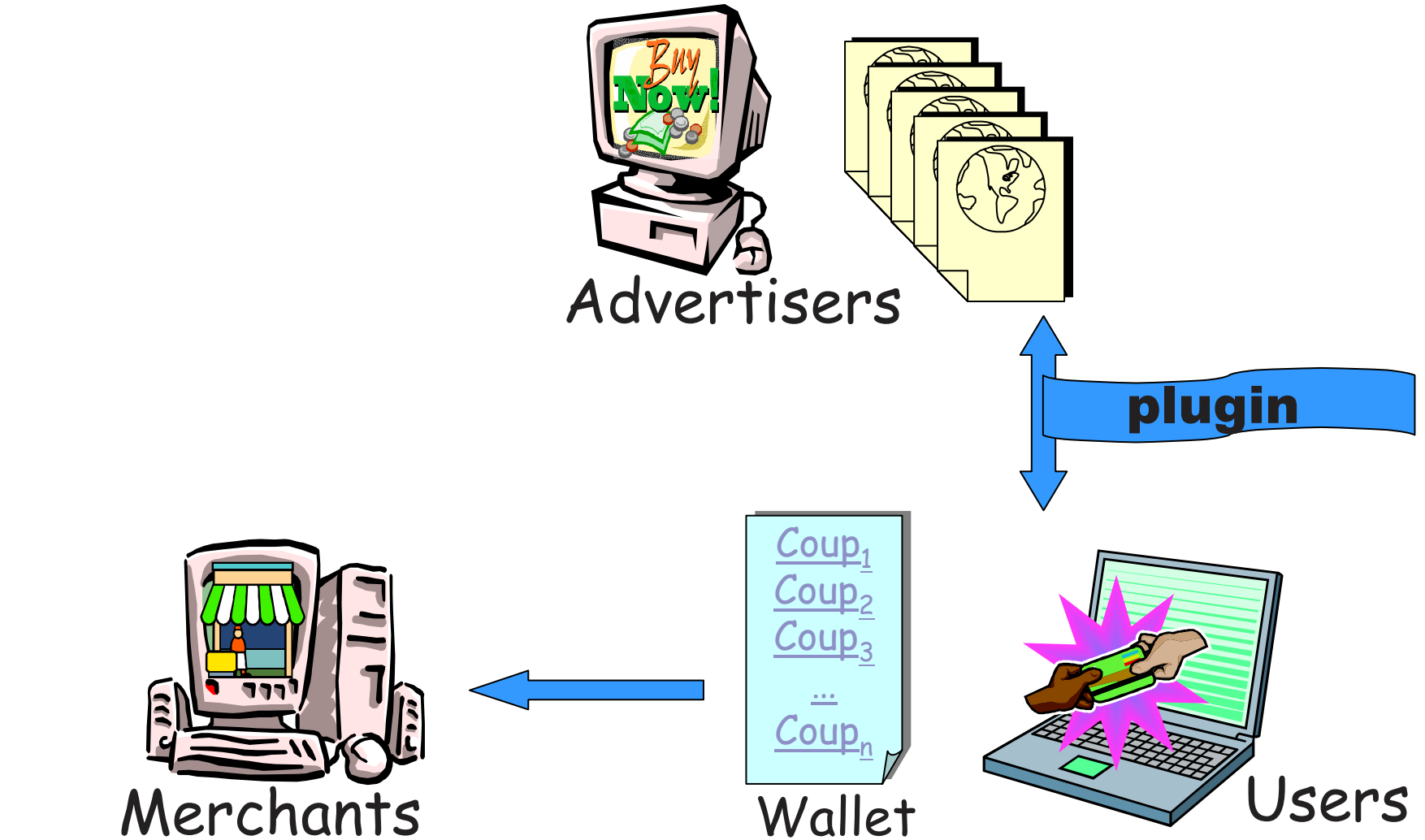
M computes $h_{KM}(m_data || q_{SN})$ and $h_{KM,A}(m_data || \dots || q_{SN})$
and accepts e-coupon **iff**

- computed values = values stored in e-coupon
- SN has not been seen before
- q_{SN} in e-coupon = q_{SN} in dependence chain (when dependence chain is available)

The extended model



Implementation



Conclusions

- Our protocol is
 - *secure:*
 - no unauthorized issuance
 - no manipulation
 - no double spending
 - *lightweight.*
 - no need of public key infrastructure
 - *practically implementable:*
 - it exploits HTTP protocol
 - *respects user's privacy.*
 - no registration required